

COMS W3261 - LECTURE 1

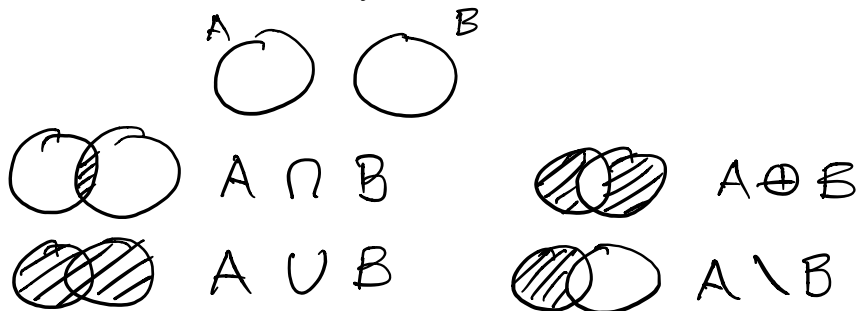
SIPSER - INTRO TO THE THEORY OF COMPUTATION, 3rd ED.

Chapter 0 - 0.2

Set: bag of unique objects

$$A = \{1, 2, 3\} \quad B = \{a, b, c\}$$

multiset: not unique. $\{1, 1, 2\}$



$$a \in B$$

$$C \subseteq B$$

$$a \notin A$$

$$C \subset B$$

$$|A| = 3,$$

Sequence := ordered set.

$$C := (a, b, c)$$

$$E := (a, a, b)$$

$$D := (b, a, c)$$

$$F := (1, 2, 4, 8, 16, \dots)$$

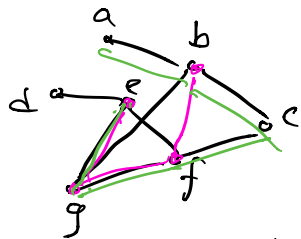
finite? Call it a k-tuple.

X - operator: Cartesian product.

$$\{a, b\} \times \{1, 2\} = \{(a, 1), (a, 2), (b, 1), (b, 2)\}$$

$$\mathbb{Z}^2, \mathbb{Z} \times \mathbb{Z} = \{ () \dots \}$$

GRAPHS:



$$G := (V, E)$$

set of
vertices

set of edges

$$V = \{a \dots g\}$$

$$E = \{ \underline{(a,b)}, \underline{(b,c)}, \underline{(c,f)} \dots \}$$

$\deg(v) :=$ number of incident edges

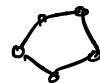
$$\deg(e) = 3$$

Subgraph := graph where, $H = (V_2, E_2)$

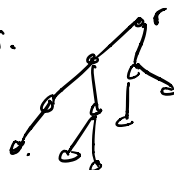
if $V_2 \subseteq V, E_2 \subseteq E$

path := (e, g, f, c, b, a)

cycle := path that ends where it starts.



tree := connected graph with no cycles.



directed graph (digraph):



(b,a)

STRINGS & LANGUAGES

alphabet: set of characters/symbols.

$$\Sigma = \{0, 1\}$$

$$= \{a, b, c \dots z\} \quad \tilde{n} \notin \tilde{\Sigma}$$

string: finite sequence of symbols.

010010 tim

Language: set of strings.

'English' := { all words w : w is in the English language }

$L :=$ { all strings $w \in \{0,1\}^n$: w is a palindrome }

$n=3$ 000
 010
 100

↑

Boolean Logic.

True = T = 1

False = F = 0

\wedge - and
 \vee - or
 \neg - not
 \rightarrow - implies.
 \otimes - xor

$$1 \wedge 1 = 1$$

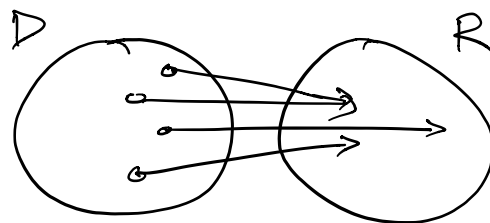
$$1 \oplus 0 = 1$$

$$1 \oplus 1 = 0$$

$$((1 \vee 0) \vee 1) \rightarrow 1$$

functions

$f: D \rightarrow R$



Proofs:

Construction —

Contradiction —

Induction —
