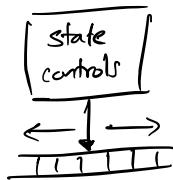


COMS W3261 - Lecture 9, Part 1:
Recognizing, Deciding, & Enumerating.

Teaser: Is the language of palindromes over the alphabet $\{0, 1\}$ Turing-recognizable?



$$L = \{ww^R, w0w^R, w1w^R \mid w \in \{0, 1\}^*\}.$$

Implementation-level description of a TM M , for L :

M_1 = "On an input s , we

- (0) accept if the tape contains 0 or 1 symbols
(shuttle back and forth to read string length.)
- (1) "Remember" the leftmost symbol and erase it.
- (2) Traverse the tape and check if the rightmost symbol matches the leftmost symbol.
 - If yes - erase the rightmost, go back to the leftmost unerased symbol, and repeat from step 0.
 - If no - reject."

✓
0000000
0101010
X

Announcements: HW #5 due 8/2/21 @ 11:59 PM EST.

See Ed for information about the final (pinned, ask questions)

Readings: Sipser 3.1 (TMs)

Sipser 3.2 (Variations on TMs, Multitape, Nondeterministic)

Sipser 3.3 ('From Turing Machines to Algorithms') Encoders.

- Today:
1. Review of TMs
 2. Variant TMs
 3. TMs \rightarrow more general notions of algorithm

1. Some TM examples (Implementation Level.)

Example.

Goal: Deciding $C = \{a^i b^j c^k \mid i, j, k \geq 1\}$.

Deciding: YES if in the language and NO otherwise

(Recognizing: YES if in the language, may loop or reject otherwise)

M_3 = "On input string w :

1. Scan left to right to ensure we have a string matching $a^+b^+c^+$.

(Reject if not.)

2. Return to the leftmost square:

3. Cross off the first a . Then, shuttle back and forth between b 's and c 's, crossing off one c for each b crossed off.
(Reject if we run out of c 's.)

4. Restore (uncross) all the b 's and return to step 2.

5. Once all a 's are crossed, accept if no c 's remain uncrossed.

Example:

~~a a a b b c c c c c~~

~~a a a b b c c c c c~~

~~||||| ||||| |||~~

(Multiplication on TMs).

Example 2. (Element Distinctness).

Goal: Decide $E = \{\#x_1 \# x_2 \# x_3 \dots \# x_e \mid \text{each } x_i \in \{0, 1\}^*\text{ and } x_i \neq x_j \text{ for } i \neq j\}$

Idea: Compare x_1 with all x_i , $i > 1$, reject if we find a match. Then compare x_2 with all x_i , $i > 2$, and so on.

M_4 = "On input w :

1. Check to make sure the string is in the right format, reject if not.

2. Accept if there is ≤ 2 inputs in the right format.

3. Otherwise, mark the first two $\#$ like this:

$\overset{\bullet}{\#} x_1 \overset{\bullet}{\#} x_2 \quad$ this is OK because we can have
 $\#, \# \in \Gamma$.

4. Scan the two strings to the right of the marked $\#$ and reject if they match.

5. If possible, move the right mark to the right and repeat step 4.

$\overset{\bullet}{\#} x_1 \overset{\bullet}{\#} x_2 \# \rightarrow \overset{\bullet}{\#} x_1 \overset{\bullet}{\#} x_2 \overset{\bullet}{\#} x_3 \dots$

If not possible (right mark is on the last $\#$), move the left mark forward and the right mark back to the next $\#$ after the left mark.

$\overset{\bullet}{\#} \dots \overset{\bullet}{\#} x_p \Rightarrow \overset{\bullet}{\#} x_1 \overset{\bullet}{\#} x_2 \overset{\bullet}{\#} \dots \#$

Finally, if we've marked the last two $\#$'s, accept.

Now: know how to nest two loops

know how to multiply

know how to check string matching and distinctness,

...

Next: discuss variants on TMs.